SA305 – Linear Programming Asst. Prof. Nelson Uhan

# Lesson 20. Geometry and Algebra of "Corner Points"

#### 0 Warm up

Example 1. Consider the system of equations

$$3x_1 + x_2 - 7x_3 = 17$$

$$x_1 + 5x_2 = 1$$

$$-2x_1 + 11x_3 = -24$$
(\*)

Let

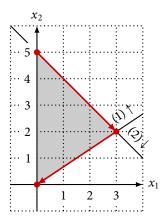
$$A = \begin{pmatrix} 3 & 1 & -7 \\ 1 & 5 & 0 \\ -2 & 0 & 11 \end{pmatrix}$$

We have that det(A) = 84.

- Does (\*) have a unique solution, no solutions, or an infinite number of solutions?
- Are the row vectors of *A* linearly independent? How about the column vectors of *A*?
- What is the rank of *A*? Does *A* have full row rank?

#### 1 Overview

- Due to convexity, local optimal solutions of LPs are global optimal solutions
  - $\Rightarrow$  Improving search finds global optimal solutions of LPs
- Last time: improving search among "corner points" of the feasible region of an LP
- Today: how can we describe "corner points" of the feasible region of an LP?
- Coming next: for LPs, is there always an optimal solution that is a "corner point"?

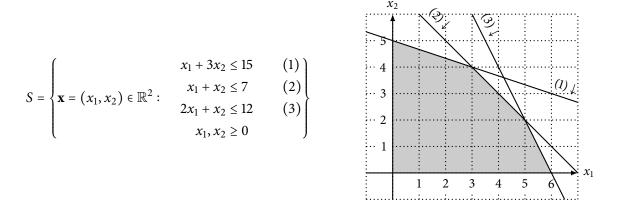


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#### 2 Polyhedra and extreme points

- A **polyhedron** is a set of vectors **x** that satisfy a finite collection of linear constraints (equalities and inequalities)
  - Also referred to as a polyhedral set
- In particular:
- Recall: the feasible region of an LP a polyhedron is a convex feasible region
- Given a convex feasible region *S*, a solution **x** ∈ *S* is an **extreme point** if there does <u>not</u> exist two distinct solutions **y**, **z** ∈ *S* such that **x** is on the line segment joining **y** and **z** 
  - i.e. there does not exist  $\lambda$  ∈ (0,1) such that  $\mathbf{x} = \lambda \mathbf{y} + (1 \lambda)\mathbf{z}$

**Example 2.** Consider the polyhedron *S* and its graph below. What are the extreme points of *S*?



• "Corner points" of the feasible region of an LP  $\Leftrightarrow$  extreme points

#### 3 Basic solutions

- In Example 2, the polyhedron is described with 2 decision variables
- Each corner point / extreme point is the intersection of 2 lines
- Equivalently, each corner point / extreme point is active at 2 distinct constraints
- Is there a connection between the number of decision variables and the number of active constraints at a corner point / extreme point?
- Convention: all variables are on the LHS of constraints, all constants are on the RHS
- A collection of constraints defining a polyhedron are **linearly independent** if the LHS coefficient matrix these constraints has full row rank

Example 3. Consider the polyhedron S given in Example 2. Are constraints (1) and (3) linearly independent?

- Given a polyhedron *S* with *n* decision variables, **x** is a **basic solution** if
  - (a) it satisfies all equality constraints
  - (b) at least n constraints are active at  $\mathbf{x}$  and are linearly independent
- **x** is a **basic feasible solution (BFS)** if it is a basic solution and satisfies all constraints of *S*

**Example 4.** Consider the polyhedron *S* given in Example 2. Verify that (3, 4) and (21/5, 18/5) are basic solutions. Are these also basic feasible solutions?

### 4 Equivalence of extreme points and basic feasible solutions

• From our examples, it appears that for polyhedra, extreme points are the same as basic feasible solutions

**Big Theorem.** Suppose *S* is a polyhedron. Then **x** is an extreme point of *S* if and only if **x** is a basic feasible solution.

- See Rader p. 243 for a proof
- We use "extreme point" and "basic feasible solution" interchangeably

## 5 Food for thought

**Problem 1.** Does a polyhedron always have an extreme point? *Hint*. Consider the following polyhedron in  $\mathbb{R}^2$ :  $S = \{(x_1, x_2) : x_1 + x_2 \ge 1\}$ .

Problem 2 (also on today's homework). Determine the extreme points of the following convex region:

$$-x_{1} + x_{2} \le 4$$
$$x_{1} - x_{2} \le 10$$
$$x_{1} + x_{2} \le 12$$
$$x_{1}, x_{2} \ge 0$$