SA305 – Linear Programming Spring 2013
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Lesson 29. Bounds and the Dual LP

1 Overview

- It is often useful to quickly generate lower and upper bounds on the optimal value of an LP
- Many algorithms for optimization problems that consider LP "subproblems" rely on this
- How can we do this?

2 Finding lower bounds

Example 1. Consider the following LP:

$$
\begin{aligned}
z^* = \text{maximize} \quad & 2x_1 + 3x_2 + 4x_3 \\
\text{subject to} \quad & 3x_1 + 2x_2 + 5x_3 \le 18 && (1) \\
& 5x_1 + 4x_2 + 3x_3 \le 16 && (2) \\
& x_1, x_2, x_3 \ge 0 && (3)
\end{aligned}
$$

We denote the optimal value of this LP by z^*. Give a feasible solution to this LP and its value. How does this value compare to z^*?

- **For a maximization LP, any feasible solution gives a lower bound on the optimal value**
- We want the greatest lower bound possible (i.e. the lower bound closest to the optimal value)

3 Finding upper bounds

- We want the lowest upper bound possible (i.e. the upper bound closest to the optimal value)
- For the LP in Example 1, we can show that the optimal value z^* is at most 27
  - Any feasible solution (x_1, x_2, x_3) must satisfy constraint (1)
  - ⇒ Any feasible solution (x_1, x_2, x_3) must also satisfy constraint (1) multiplied by 3/2 on both sides:

- The nonnegativity bounds (3) imply that any feasible solution (x_1, x_2, x_3) must satisfy

- Therefore, any feasible solution, including the optimal solution, must have value at most 27

- We can do better: we can show $z^* \leq 25$:
  - Any feasible solution (x_1, x_2, x_3) must satisfy constraints (1) and (2)
  - $\Rightarrow$ Any feasible solution (x_1, x_2, x_3) must also satisfy $\left(\frac{1}{2} \times \text{constraint (1)}\right) + \text{constraint (2)}$:

  - The nonnegativity bounds (3) then imply that any feasible solution (x_1, x_2, x_3) must satisfy

Example 2. Combine the constraints (1) and (2) of the LP in Example 1 to find a better upper bound on z^* than 25.

- Let's generalize this process of combining constraints
- Let y_1 be the "multiplier" for constraint (1), and let y_2 be the "multiplier" for constraint (2)
- We require $y_1 \geq 0$ and $y_2 \geq 0$ so that multiplying constraints (1) and (2) by these values keeps the inequalities as "$\leq$"
- We also want:

- Since we want the lowest upper bound, we want:

- Putting this all together, we can find the multipliers that find the best lower upper bound with the following LP!

$$\begin{aligned} \text{minimize} \quad & 18y_1 + 16y_2 \\ \text{subject to} \quad & 3y_1 + 5y_2 \geq 2 \\ & 2y_1 + 4y_2 \geq 3 \\ & 5y_1 + 3y_2 \geq 4 \\ & y_1 \geq 0, y_2 \geq 0 \end{aligned}$$

  - This is the **dual LP**, or simply the **dual** of the LP in Example 1
  - The LP in example is referred to as the **primal LP** or the **primal** – the original LP

4 In general...

- Every LP has a dual
- For minimization LPs
  - Any feasible solution gives an upper bound on the optimal value
  - One can construct a dual LP to give the greatest lower bound possible
- We can generalize the process we just went through to develop some mechanical rules to construct duals

5 Constructing the dual LP

0. Rewrite the primal so all variables are on the LHS and all constants are on the RHS
1. Assign each primal constraint a corresponding **dual variable** (multiplier)
2. Write the dual objective function
   - The objective function coefficient of a dual variable is the RHS coefficient of its corresponding primal constraint
   - The dual objective sense is the opposite of the primal objective sense
3. Write the dual constraint corresponding to each primal variable
   - The dual constraint LHS is found by looking at the coefficients of the corresponding primal variable ("go down the column")
   - The dual constraint RHS is the objective function coefficient of the corresponding primal variable
4. Use the **SOB rule** to determine dual variable bounds (≥ 0, ≤ 0, free) and dual constraint comparisons ($\leq$, $\geq$, $=$)

	max LP	$\leftrightarrow$	min LP	
sensible	$\leq$ constraint	$\leftrightarrow$	$y_i \geq 0$	sensible
odd	$=$ constraint	$\leftrightarrow$	y_i free	odd
bizarre	$\geq$ constraint	$\leftrightarrow$	$y_i \leq 0$	bizarre
sensible	$x_i \geq 0$	$\leftrightarrow$	$\geq$ constraint	sensible
odd	x_i free	$\leftrightarrow$	$=$ constraint	odd
bizarre	$x_i \leq 0$	$\leftrightarrow$	$\leq$ constraint	bizarre

Example 3. Take the dual of the following LP:

$$\begin{aligned} \text{minimize} \quad & 10x_1 + 9x_2 - 6x_3 \\ \text{subject to} \quad & 2x_1 - x_2 \geq 3 \\ & 5x_1 + 3x_2 - x_3 \leq 14 \\ & x_2 + x_3 = 1 \\ & x_1 \geq 0, x_2 \leq 0, x_3 \geq 0 \end{aligned}$$

- **The dual of the dual is the primal**
  - Try it with the dual you just found