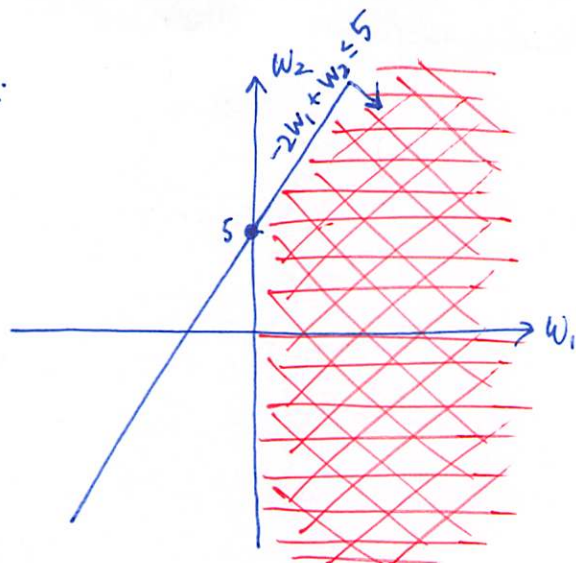


P1.



Graph of feasible region ↷

a. minimize  $3w_1 - w_2$   
⇒ opt. soln. is  $(0, 5)$

b. minimize  $w_1$   
⇒ opt. soln. is the intersection of the  $w_2$ -axis and the feasible region

c. minimize  $-w_1$   
⇒ model is unbounded (obj. fn. contours with better obj. fn. value "move to the right" without limit.)

P2. a. Current basis =  $\{x_1, x_3\} \Rightarrow x_2 = x_4 = 0$  at current BFS:

$$\left. \begin{array}{l} -x_1 + 4x_3 = 13 \\ 2x_1 = 2 \end{array} \right\} \Rightarrow x_1 = 1, x_3 = \frac{14}{4} = \frac{7}{2}$$

$\Rightarrow$  current BFS =  $(1, 0, \frac{7}{2}, 0)$ .

b.  $\vec{d}^{x_2} = (d_{x_1}, 1, d_{x_3}, 0)$

$$-d_{x_1} + 4d_{x_3} = -1$$

$$2d_{x_1} = -6$$

$$\Rightarrow d_{x_1} = -3, d_{x_3} = -1$$

$$\Rightarrow \vec{d}^{x_2} = (-3, 1, -1, 0)$$

$$\vec{d}^{x_4} = (d_{x_1}, 0, d_{x_3}, 1)$$

$$-d_{x_1} + 4d_{x_3} = -21$$

$$2d_{x_1} = 2$$

$$\Rightarrow d_{x_1} = 1, d_{x_3} = -5$$

$$\Rightarrow \vec{d}^{x_4} = (1, 0, -5, 1)$$

c.  $\bar{c}_{x_2} = -30 + 1 = -29$

$\Rightarrow \vec{d}^{x_2}$  not improving

$$\bar{c}_{x_4} = 10$$

$\Rightarrow \vec{d}^{x_4}$  improving.

d. Choose  $\vec{d}^{x_4}$ :  $x_4$  is the entering variable

$$\lambda_{\max} = \min \left\{ \frac{7/2}{5} \right\} = \frac{7}{10} \quad x_3 \text{ is the leaving variable.}$$

$$\Rightarrow \text{new BFS} = (1, 0, \frac{7}{2}, 0) + \frac{7}{10}(1, 0, -5, 1) = (\frac{17}{10}, 0, 0, \frac{7}{10})$$

$$\text{new basis} = \{x_1, x_4\}.$$

P3. a.  $\min 5y_1 + 8y_2$

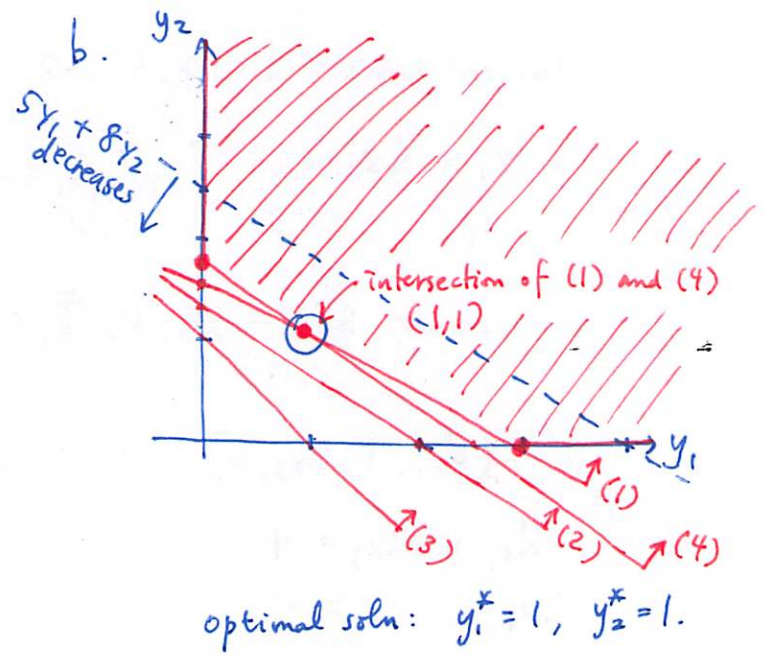
s.t.  $y_1 + 2y_2 \geq 3$  (1)

$2y_1 + 3y_2 \geq 4$  (2)

$y_1 + y_2 \geq 1$  (3)

$2y_1 + 3y_2 \geq 5$  (4)

$y_1, y_2 \geq 0.$



c. Dual complementary slackness: dual constraint not active  
 $\Rightarrow$  corresponding primal variable = 0.

From part b, it is clear that constraints (2) and (3) are not active at  $(1, 1)$ .  $\Rightarrow$  In an optimal solution to [P], we must have  $x_2 = x_3 = 0$ .

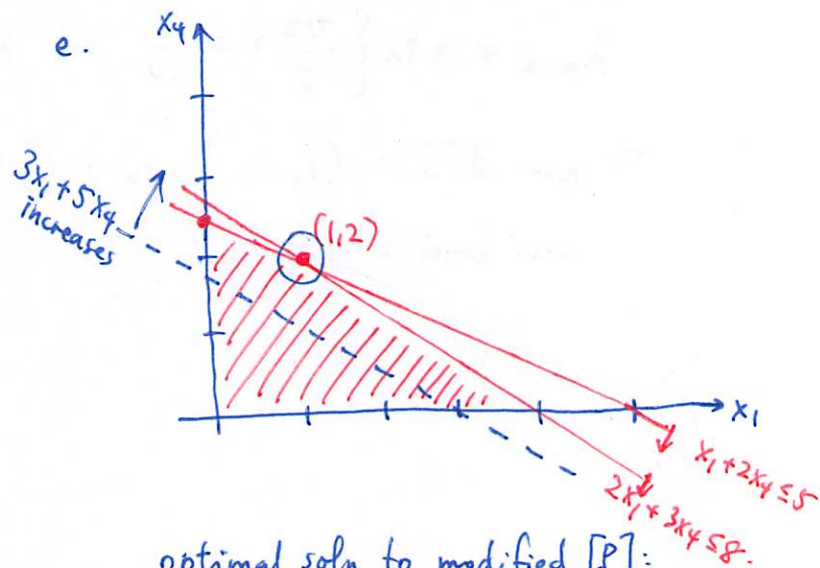
d. modified [P]:

$\max 3x_1 + 5x_4$

s.t.  $x_1 + 2x_4 \leq 5$

$2x_1 + 3x_4 \leq 8$

$x_1, x_4 \geq 0$

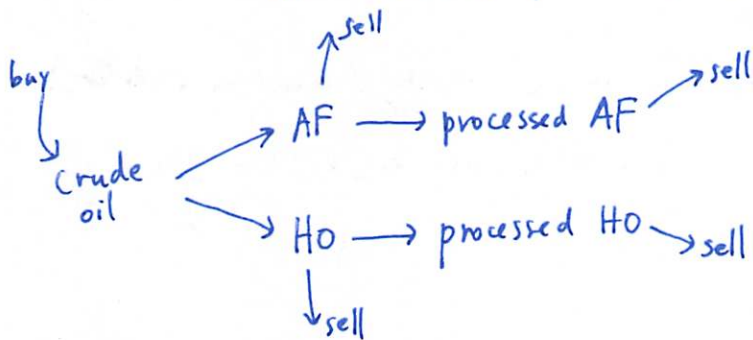


optimal soln. to [P]:  $(1, 0, 0, 2)$

P4. Symbolic input parameters:

- $c$  = cost of crude oil per 1000 barrels
- $p_a$  = price of unprocessed aviation fuel per 1000 barrels
- $p_h$  = price of unprocessed heating oil per 1000 barrels
- $q_a$  = price of processed aviation fuel per 1000 barrels
- $q_h$  = price of processed heating oil per 1000 barrels
- $t_a$  = time to process aviation fuel per 1000 barrels
- $t_h$  = time to process heating oil per 1000 barrels

$B$  = available crude oil, in 1000s.



Decision variables: (in 1000 barrels)

- $z$  = amt. of crude oil to buy
- $x_a$  = amt. of unprocessed AF to sell
- $x_h$  = amt. of unprocessed HO to sell
- $y_a$  = amt. of processed AF to sell
- $y_h$  = amt. of processed HO to sell.

Model:

$$\max \quad p_a x_a + p_h x_h + q_a y_a + q_h y_h - cz \quad (\text{total profit})$$

$$\text{s.t.} \quad \frac{3}{4}z = x_a + y_a \quad (\text{crude oil} \rightarrow \text{AF})$$

$$\frac{1}{4}z = x_h + y_h \quad (\text{crude oil} \rightarrow \text{HO})$$

$$t_a y_a + t_h y_h \leq T \quad (\text{cracker time})$$

$$z \leq B \quad (\text{available crude oil})$$

$$x_a, y_a, x_h, y_h, z \geq 0.$$

PS. The objective fn. vector  $\vec{c} = (3, 11, -8, 0)$

a.  $\vec{d}^{w_4} = (1, 0, -4, 1)$  does NOT lead to a conclusion that the LP is unbounded, since its components are not all nonnegative.

b.  $\vec{d}^{w_4} = (1, 3, 0, 1)$  has associated reduced cost  $\bar{c}_{w_4} = 36$ .  
Since the LP is minimizing,  $\vec{d}^{w_4}$  is not improving. So, even though all components of  $\vec{d}^{w_4}$  are nonnegative, we cannot conclude that the LP is unbounded.

c.  $\vec{d}^{w_4} = (1, 0, 3, 1)$  has associated reduced cost  $\bar{c}_{w_4} = -21$ , and so  $\vec{d}^{w_4}$  is improving, with all nonnegative components. Therefore, we can conclude that the LP is unbounded.

d.  $\vec{d}^{w_4} = (-1, 1, -2, 1)$ . Similar to part a.

P6. Symbolic input parameters:

$P$  = set of presents

$C$  = set of children

$V_{ij}$  = happiness of child  $i$  w/present  $j$  for  $i \in C, j \in P$

$b_j$  = # present  $j$  available for  $j \in P$ .

Decision variables:  $x_{ij}$  = # present  $j$  given to child  $i$  for  $i \in C, j \in P$ .

Model:  $\max \min \left\{ \sum_{j \in P} V_{ij} x_{ij} : i \in C \right\}$

↑ happiness of child  $i$

s.t.  $\sum_{i \in C} x_{ij} \leq b_j$  for  $j \in P$  (available presents)

$x_{ij} \geq 0$  for  $i \in C, j \in P$ .

convert to LP



$\max z$

s.t.  $z \leq \sum_{j \in P} V_{ij} x_{ij}$  for  $i \in C$

$\sum_{i \in C} x_{ij} \leq b_j$  for  $j \in P$

$x_{ij} \geq 0$  for  $i \in C, j \in P$ .