

Lesson 15. Improving Search: Finding Better Solutions

1 A general optimization model

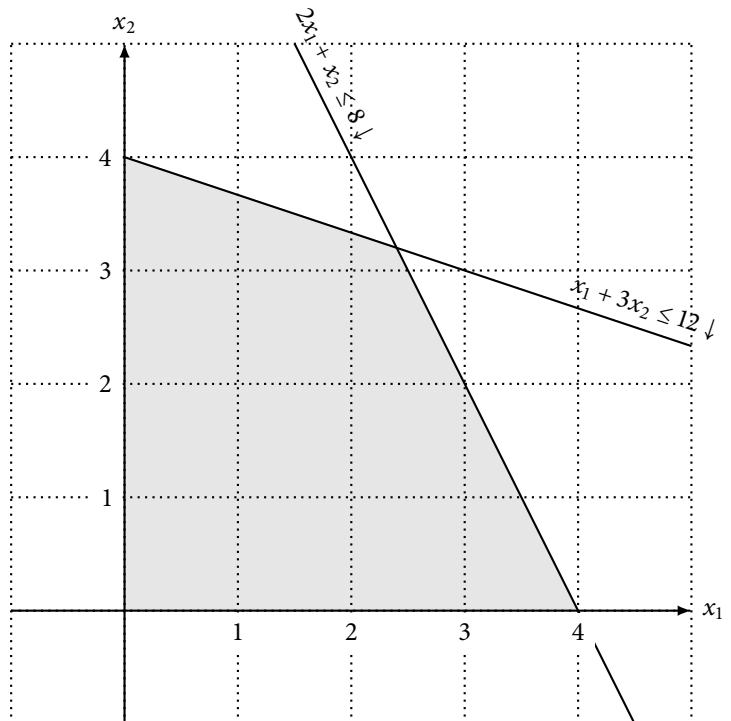
- For the next few lessons, we will consider a general optimization model
- Decision variables: x_1, \dots, x_n
 - Recall: a feasible solution to an optimization model is a choice of values for all decision variables that satisfies all constraints
- Easier to refer to a feasible solution as a vector: $\mathbf{x} = (x_1, \dots, x_n)$
- Let $f(\mathbf{x})$ and $g_i(\mathbf{x})$ for $i \in \{1, \dots, m\}$ be multivariable functions in \mathbf{x} , not necessarily linear
- Let b_i for $i \in \{1, \dots, m\}$ be constant scalars

$$\begin{aligned} &\text{minimize/maximize } f(\mathbf{x}) \\ &\text{subject to } g_i(\mathbf{x}) \begin{cases} \leq \\ \geq \\ = \end{cases} b_i \quad \text{for } i \in \{1, \dots, m\} \end{aligned} \quad (*)$$

- Linear programs fit into this framework

Example 1.

$$\begin{aligned} &\text{maximize } 4x_1 + 2x_2 \\ &\text{subject to } x_1 + 3x_2 \leq 12 \quad (1) \\ &\quad \quad \quad 2x_1 + x_2 \leq 8 \quad (2) \\ &\quad \quad \quad x_1 \geq 0 \quad (3) \\ &\quad \quad \quad x_2 \geq 0 \quad (4) \end{aligned}$$



2 Improving search algorithms, informally

- Idea:
 - Start at a feasible solution
 - Repeatedly move to a “close” feasible solution with better objective function value
- The **neighborhood** of a feasible solution is the set of all feasible solutions “close” to it
 - We can define “close” in various ways to design different types of algorithms
- Let’s start formalizing these ideas

3 Locally and globally optimal solutions

- ε -**neighborhood** $N_\varepsilon(\mathbf{x})$ of a solution $\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{R}^n$ (where $\varepsilon > 0$):

$$N_\varepsilon(\mathbf{x}) = \{\mathbf{y} \in \mathbb{R}^n : d(\mathbf{x}, \mathbf{y}) \leq \varepsilon\}$$

where $d(\mathbf{x}, \mathbf{y})$ is the distance between solution \mathbf{x} and \mathbf{y}

- A feasible solution \mathbf{x} to optimization model (*) is **locally optimal** if for some value of $\varepsilon > 0$:

$$f(\mathbf{x}) \text{ is better than } f(\mathbf{y}) \quad \text{for all feasible solutions } \mathbf{y} \in N_\varepsilon(\mathbf{x})$$

- A feasible solution \mathbf{x} to optimization model (*) is **globally optimal** if:

$$f(\mathbf{x}) \text{ is better than } f(\mathbf{y}) \quad \text{for all feasible solutions } \mathbf{y}$$

- Also known simply as an **optimal solution**

- Global optimal solutions are locally optimal, but not vice versa
- In general: harder to check for global optimality, easier to check for local optimality

4 The improving search algorithm

- 1: Find an initial feasible solution \mathbf{x}^0
- 2: Set $k = 0$
- 3: **while** \mathbf{x}^k is not locally optimal **do**
- 4: Determine a new feasible solution \mathbf{x}^{k+1} that improves the objective value at \mathbf{x}^k
- 5: Set $k = k + 1$
- 6: **end while**

- Generates sequence of feasible solutions $\mathbf{x}^0, \mathbf{x}^1, \mathbf{x}^2, \dots$
- In general, improving search converges to a local optimal solution, not a global optimal solution
- Let’s concentrate on line 4 – finding better feasible solutions

5 Moving between solutions

- How do we move from one solution to the next?

$$\mathbf{x}^{k+1} = \mathbf{x}^k + \lambda \mathbf{d}$$

- In Example 1:

6 Improving directions

- We want to choose \mathbf{d} so that \mathbf{x}^{k+1} has a better value than \mathbf{x}^k
- \mathbf{d} is an **improving direction** at solution \mathbf{x}^k if

$$f(\mathbf{x}^k + \lambda \mathbf{d}) \text{ is better than } f(\mathbf{x}^k) \text{ for all positive } \lambda \text{ "close" to } 0$$

- How do we find an improving direction?
- The **directional derivative** of f in the direction \mathbf{d} at solution \mathbf{x}^k is

- Maximizing f : \mathbf{d} is an improving direction at \mathbf{x}^k if

- Minimizing f : \mathbf{d} is an improving direction at \mathbf{x}^k if

- In Example 1:

- For linear programs in general: if \mathbf{d} is an improving direction at \mathbf{x}^k , then $f(\mathbf{x}^k + \lambda \mathbf{d})$ improves as $\lambda \rightarrow \infty$

7 Step size

- We have an improving direction \mathbf{d} – now how far do we go?
- One idea: find maximum value of λ so that $\mathbf{x}^k + \lambda\mathbf{d}$ is still feasible
- Graphically, we can eyeball this
- Algebraically, we can compute this – in Example 1:



8 Feasible directions

- Some improving directions don't lead to any new feasible solutions
- \mathbf{d} is a **feasible direction** at feasible solution \mathbf{x}^k if $\mathbf{x}^k + \lambda\mathbf{d}$ is feasible for all positive λ “close” to 0
- Again, graphically, we can eyeball this
- A constraint is **active** at feasible solution \mathbf{x} if it is satisfied with equality

- For linear programs:

- We have constraints of the form:

$$a_1x_1 + a_2x_2 + \cdots + a_nx_n \leq b$$

$$a_1x_1 + a_2x_2 + \cdots + a_nx_n \geq b$$

$$a_1x_1 + a_2x_2 + \cdots + a_nx_n = b$$

- We can rewrite these constraints using vector notation:



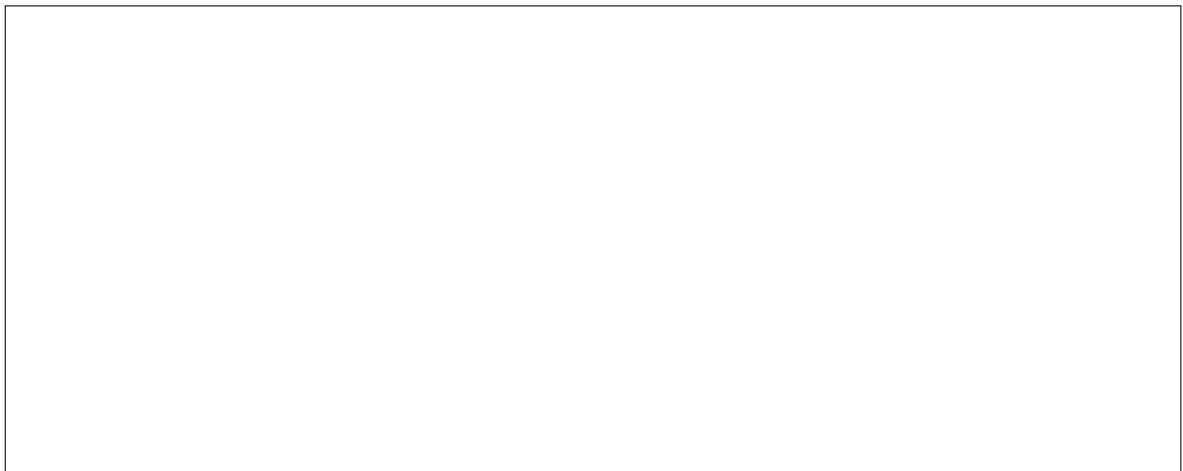
- \mathbf{d} is a feasible direction at \mathbf{x} if

- ◊ $\mathbf{a}^T \mathbf{d} \leq 0$ for each active constraint of the form $\mathbf{a}^T \mathbf{x} \leq b$

- ◊ $\mathbf{a}^T \mathbf{d} \geq 0$ for each active constraint of the form $\mathbf{a}^T \mathbf{x} \geq b$

- ◊ $\mathbf{a}^T \mathbf{d} = 0$ for each active constraint of the form $\mathbf{a}^T \mathbf{x} = b$

- In Example 1:



9 Detecting unboundedness

- Suppose \mathbf{d} is an improving direction at feasible solution \mathbf{x}^k to a linear program
- Also, suppose $\mathbf{x}^k + \lambda\mathbf{d}$ is feasible for all $\lambda \geq 0$
- What can you conclude?

10 Summary

- Line 4 boils down to finding an improving and feasible direction \mathbf{d} and an accompanying step size λ
- We discussed conditions on whether a direction is improving and feasible
- We don't know how to systematically find such directions... yet